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WHAT IS CLAIMED IS:

- 1. A method for countering unauthorized decryption comprises a step of scrambling at least one correlation between a data decryption processing in a hardware and at least one respective hardware operational phenomenon by randomly changing at least one arithmetic operation order in the data decryption processing.
- 2. The method for countering unauthorized decryption according to Claim 1, wherein the hardware operational phenomenon is power consummated by the hardware to execute the data decryption processing.
 - 3. The method for countering unauthorized decryption according to Claim 1, wherein the hardware is a IC card, a PDA, or a cellular phone.
 - 4. The method for countering unauthorized decryption according to Claim 1, wherein the data decryption processing is executed to decrypt data encrypted by a RSA encryption processing or an elliptic encryption processing.
- 5. The method for countering unauthorized decryption according to Claim 1, wherein the correlation is scrambled by an arithmetic operation method implemented by an information processing apparatus comprising the steps of:

for two integers K1 and K2, when finding a value F(K, A) of a function F satisfying F(K1+K2, A)=F(K1, A) OF(K2, A) (O denotes an arithmetic operation in a communitative semigroup S. K designates an integer and A designates an

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element of S), decomposing the K to the sum of m integers K[0] + K[1] + ... K[m-1];

using T(0), T(1), ... T(m-1) resulted from rearranging a string of integers 0, 1, ... m-1 by permutation T; and

- operating on terms F(K[T(0)], A) to F(K[T(m-1)], A) on the right side of F(K, A) = F(K[T(0)], A) OF(K[T(1)], A) O ... F(K[T(m-1)], A) ... ("expression 1") in an order of F(K[T(0)], A), F(K[T(1)], A), ... F(K[T(m-1)], A) to find F(K, A).
- 10 6. The method according to claim 5, whereby the permutation processing, the permutation T prevents predicting any post-permutation data from pre-permutation data, or the permutation T is performed based on a dummy random number, and whereby the permutation processing is performed each time the expression 1 is performed.

7. The method according to claim 5, wherein the S is a commutative semigroup in which, for a set consisting of residues by an integer N (N \geq 2), the arithmetic operation O of a modular multiplication operation AOB=A*B mod N is introduced, and the F satisfies F(K, A)=A^K mod N (A^K denotes the K-th power of A).

- 8. The method according to claim 5, wherein the information processing apparatus is installed on an IC card, a cellular phone, or a PDA.
- 25 9. The method according to claim 7, wherein the integer K is split in a form of

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$$K[j] = u^*((2^t)^j) (0 \le u \le (2^t) - 1, t = 1, 2, ...)$$

- 10. The method according to claim 9, whereby the permutation processing, the permutation T is performed based on an information source prevents predicting any post-permutation data from pre-permutation data, or the permutation T is performed based on a dummy random number, and whereby the permutation processing is performed each time the expression 1 is performed.
- 11. The method according to claim 9, wherein the integer K is split in a form of
 10 K[j] = u*((2^t)^j) (0≤u≤(2^t)-1, t = 1,2,...)
 - 12. The method according to claim 5, wherein the S is a Mordell-Weil group on an elliptic curve E defined on a finite field GF(p) (p is a prime number) or $GF(2^n)$ (n is an integer equal to or greater than 1), and an expression F(K,A)=KA is satisfied, wherein the A denotes a point on the elliptic curve E, the KA denotes the arithmetic operation O performed on K number of As such that the KA denotes A OAOA ... OA (K number) when the K is positive, or (-A)O(-A)O(-A) ... O (-A)(|K|) number) when the K is negative, and 0 (the point at infinity) on the E when the K is 0, the O denotes an addition operation in the Mordell-Weil group, and the -A is an inverse in the Mordell-Weil group of the A.
 - 13. The method according to claim 12, wherein the information processing apparatus is installed on an IC card.
- 25 14. The method according to claim 12, wherein the integer K is split in a form of

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$$K[j] = u*((2^t)^j) (0 \le u \le (2^t)-1, t = 1,2,...)$$

- 15. The method according to claim 14, wherein the information processing apparatus is installed on an IC card.
- 16. An apparatus for countering unauthorized decryption comprises means for scrambling at least one correlation between a data decryption processing in a hardware and at least one respective hardware operational phenomenon by randomly changing at least one arithmetic operation order in the data decryption processing.
- 17. The apparatus for countering unauthorized decryption according to Claim 16, wherein the hardware operational phenomenon is power consummated by the hardware to execute the data decryption processing.
- 18. The apparatus for countering unauthorized decryption according to Claim 16, wherein the hardware is a IC card, a PDA, or a cellular phone.
- 19. A software product for countering unauthorized decryption comprises a module for scrambling at least one correlation between a data decryption processing in a hardware and at least one respective hardware operational phenomenon by randomly changing at least one arithmetic operation order in the data decryption processing.
- 25 20. The software product for countering unauthorized decryption according to

Claim 19, wherein the hardware operational phenomenon is power consummated by the hardware to execute the data decryption processing.